

Dependency Parsing with Bounded Block Degree and Well-nestedness via Lagrangian Relaxation and Branch-and-Bound

Caio Corro, Joseph Le Roux, Mathieu Lacroix, Antoine Rozenknop and Roberto Wolfler-Calvo

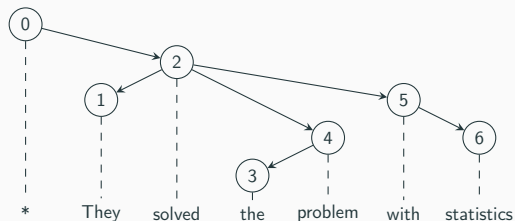
August 7-12

Université Paris 13 – LIPN

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Dependency trees

- Association of each word of sentence with a vertex
- Dependency tree: spanning tree rooted at 0

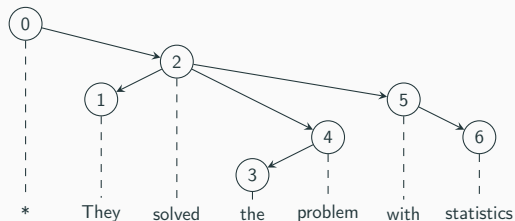


Dependency parsing

- Set of valid dependency trees for sentence x : Y_x
- Arc factored model: $score(y) = \sum_{a \in Y} score(a)$
- Dependency parsing: $\hat{y}_x = \arg \max_{y \in Y_x} score(y)$

Dependency trees

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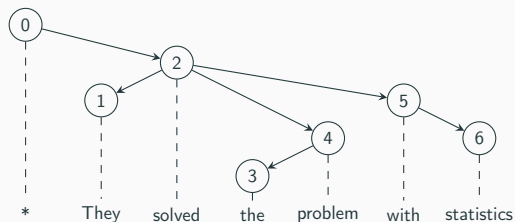


Structural properties [Bodirsky et al. 2009; Kuhlmann 2010]

Non-projective \longleftrightarrow **Projective**

Dependency trees

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Structural properties [Bodirsky et al. 2009; Kuhlmann 2010]



Distribution of dependency tree characteristics

	English (PTB/LTH)		German (SPMRL)		Dutch (UD)	
	WN	WN	WN	WN	WN	WN
BD 1	92.26		67.60		69.13	
BD 2	7.58	0.12	27.12	0.79	28.50	0.08
BD 3	0.12	0.01	3.86	0.30	2.24	0.01
BD 4	0.00	0.00	0.19	<0.01	0.04	0.00
BD > 4	0.00	0.00	0.11	<0.01	0.00	0.00
	Spanish (UD)		Portuguese (UD)			
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- **Blue** + **Purple**: $\approx 99\%$ of the dependency trees

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- **Blue**: Projective dependency trees
- **Blue** + **Purple**: $\approx 99\%$ of the dependency trees
- **Blue** + **Purple** + **Red**: Non-projective dependency trees

Observation

- Projective parsing: does not correctly cover datasets
- Non-projective parsing: produce invalid structures

Problem

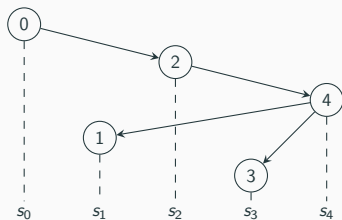
- WN and k-BBD parsing: no tractable algorithm

Contribution

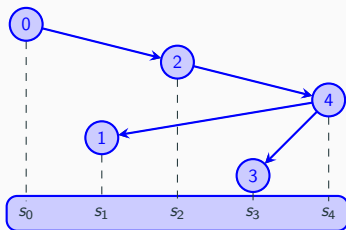
- First efficient parsing algorithm based on Lagrangian Relaxation

1. Introduction
2. Dependency tree characterization
3. Existing parsing algorithms
4. Novel characterization based on arc-sets
5. Efficient parsing with fine-grained constraints
6. Experiments
7. Conclusion

Yield of a node v : set of all nodes reachable from v

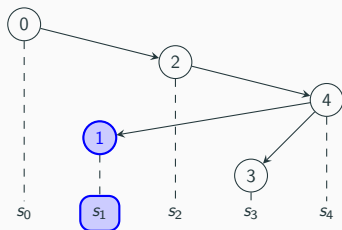


Yield of a node v : set of all nodes reachable from v



$$Yield(0) = \{0, 1, 2, 3, 4\}$$

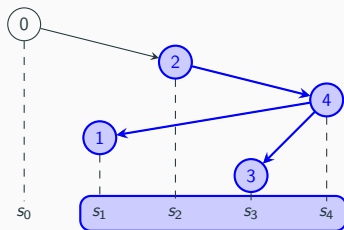
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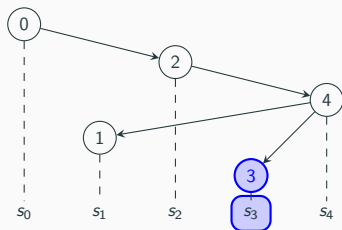


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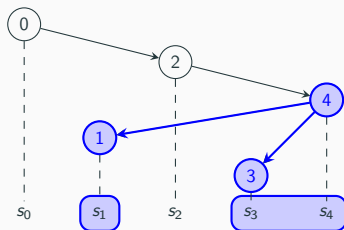
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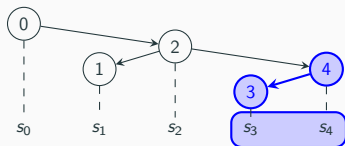
$$Yield(3) = \{3\}$$

$$Yield(4) = \{3, 4\}$$

Structural properties of dependencies

Projective dependency trees

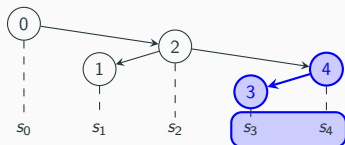
⇒ Trees with contiguous yields only



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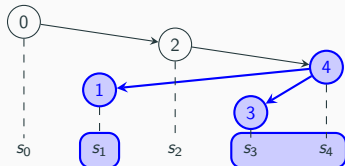
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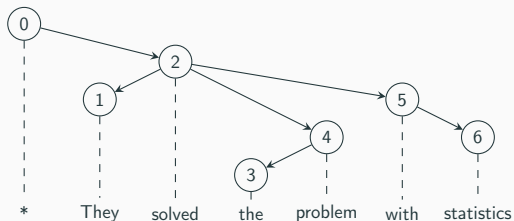
Non-projective dependency trees

⇒ Unconstrained trees

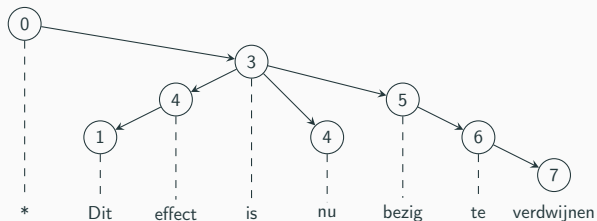


Example: Projective dependency trees

- English

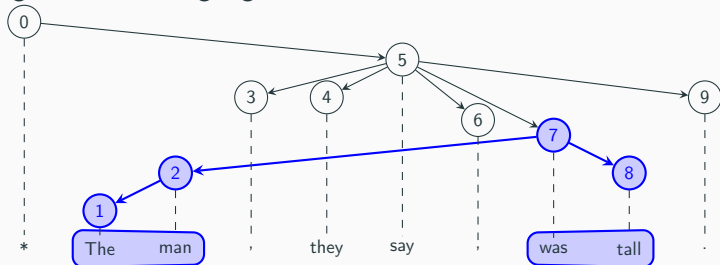


- Dutch

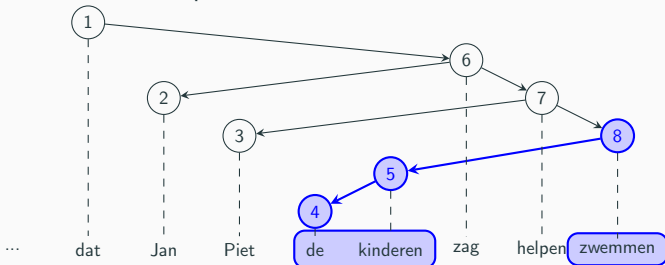


Example: Non-projective dependency trees

- English: surrounding argument



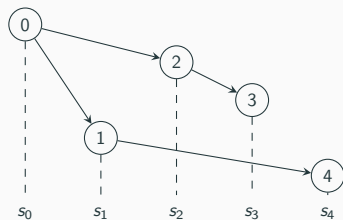
- Dutch: cross-serial dependencies



Structural properties (1/2): k-BBD

k-Bounded Block Degree (k-BBD)

- BD of a vertex: number of contiguous intervals described by its yield
- BD of a tree: the maximal block degree of its vertices
- k-BBD tree: tree with a BD less or equal to k

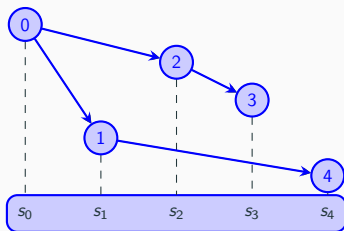


Tree of block degree 2

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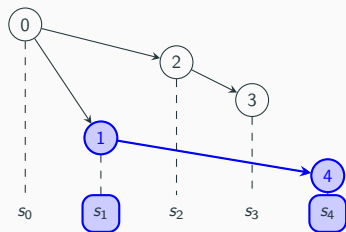
$BD(0) = 1$

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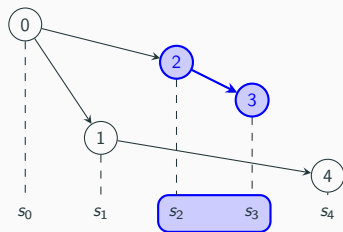
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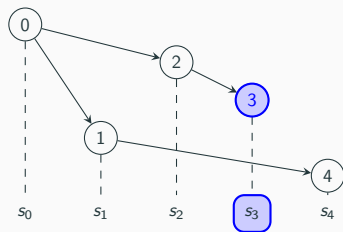
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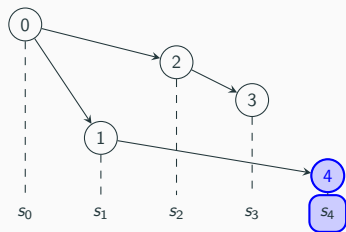
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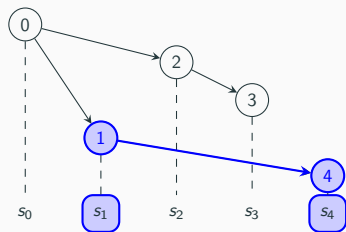
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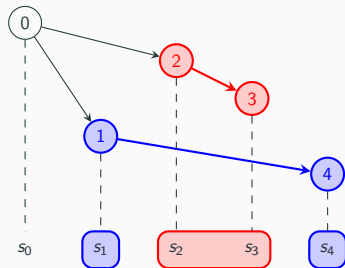
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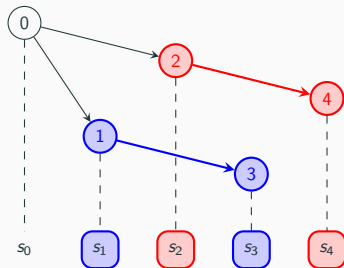
Structural properties (2/2): WN

Well-nestedness (WN)

- Interleaving sets I_1, I_2 : there exist $i, j \in I_1$ and $k, l \in I_2$ such that $i < k < j < l$
- *Well-nested* tree: does not contain two vertices whose yields are disjoint and interleave



Well-nested tree



Not well-nested tree

Complexity (arc-factored)

Non-projective	$O(n^2)$	[McDonald et al. 2005]
Projective	$O(n^3)$	[Eisner 2000]
WN + 2-BBD	$O(n^7)$	[Gómez-Rodríguez et al. 2009]
WN + k-BBD, $k \geq 2$	$O(n^{5+2(k-1)})$	[Gómez-Rodríguez et al. 2009]

Remark

Projective \Leftrightarrow 1-BBD and WN

Tractability

- Non-projective and projective: tractable
- WN + k-BBD: not tractable

Non-projective dependency parsing

Integer Linear Program for non-projective parsing

$z \in R^A$: incidence vector such that $z_a = 1$ iff arc a is in the tree.

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$$\text{s.t.} \quad \sum_{a \in \delta^{\text{in}}(v)} z_a = 1 \quad \forall v \in V^+ \quad \text{One head/word} \quad (2)$$

$$\sum_{a \in \delta^{\text{in}}(W)} z_a \geq 1 \quad \forall W \subseteq V^+ \quad \text{Connectedness} \quad (3)$$

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Efficient decoding

In practice: (directed) *Maximum Spanning Tree* (MST) algorithm
[Schrijver 2003; McDonald et al. 2005]

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Problem enhancement

\Rightarrow Integrating fine-grained structural constraints ?

k-Bounded Block Degree Constraint

Definition

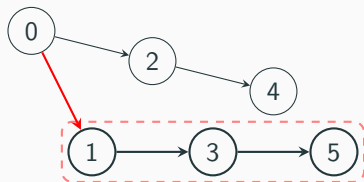
\mathcal{W}^{k+1} : vertex subsets describing at least $k + 1$ non-adjacent intervals

k-Bounded Block Degree Constraint

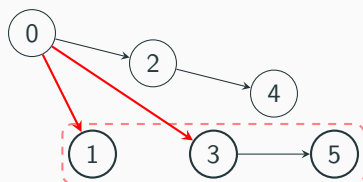
Definition

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Example with $k = 2$ and $\{1, 3, 5\} \in \mathcal{W}^3$



Not 2-BBD



2-BBD

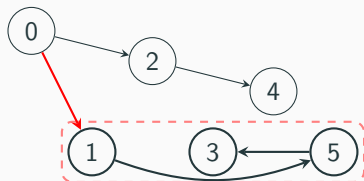
\rightarrow : arcs adjacent to the vertex subset $\{1, 3, 5\}$

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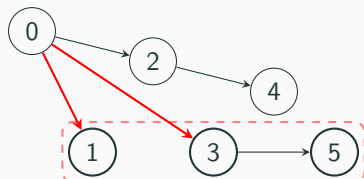
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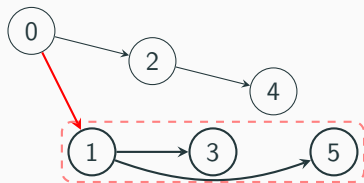
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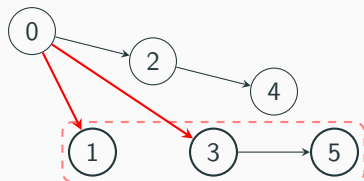
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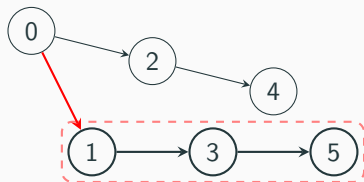
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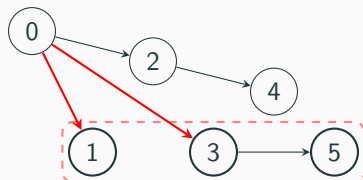
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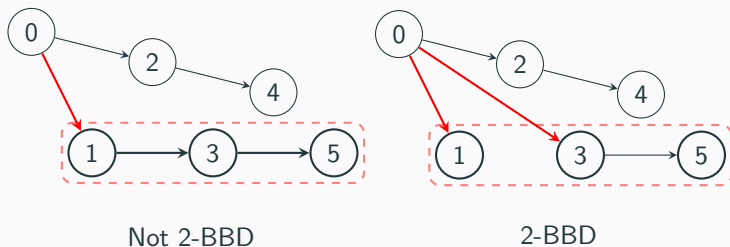
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k-Bounded Block Degree Constraint

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Example with $k = 2$ and $\{1, 3, 5\} \in \mathcal{W}^3$



\rightarrow : arcs adjacent to the vertex subset $\{1, 3, 5\}$

Constraint

For each vertex subset $W \in \mathcal{W}^{\geq k+1} \Rightarrow$ At least two adjacent arcs

Well-nestedness constraint

Definition

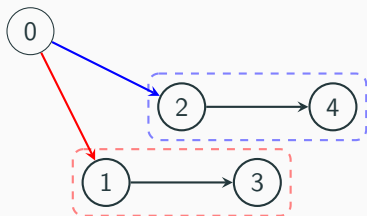
\mathcal{I} : family of couples of disjoint interleaving vertex subsets

Well-nestedness constraint

Definition

\mathcal{I} : family of couples of disjoint interleaving vertex subsets

Example with $(\{1, 3\}, \{2, 4\}) \in \mathcal{I}$

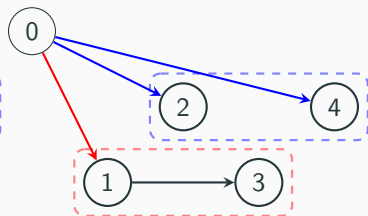


Not Well-nested

$$Yield(1) = \{1, 3\}$$

$$Yield(2) = \{2, 4\}$$

$$1 < 2 < 3 < 4$$



Well-nested

$$Yield(1) = \{1, 3\}$$

$$Yield(2) = \{2\}$$

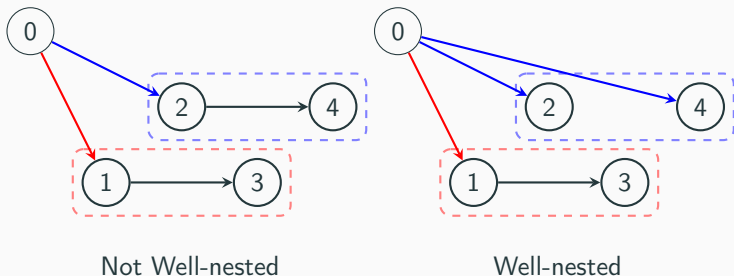
$$Yield(4) = \{4\}$$

Well-nestedness constraint

Definition

\mathcal{I} : family of couples of disjoint interleaving vertex subsets

Example with $(\{1, 3\}, \{2, 4\}) \in \mathcal{I}$



Constraint

For each couple $(l_1, l_2) \in \mathcal{I} \Rightarrow$ At least two adjacent arcs for l_1 or l_2

Full ILP: parsing with k-BBD and WN constraints

$$\max_z \sum_{a \in A} \text{score}(a) \times z_a \quad \text{Arc-factored} \quad (5)$$

$$\text{s.t. } z \in Z \quad \text{Non-projective} \quad (6)$$

$$\sum_{a \in \delta(W)} z_a \geq 2 \quad \forall W \in \mathcal{W}^{\geq k+1} \quad k\text{-BBD} \quad (7)$$

$$\sum_{a \in \delta(l_1)} z_a + \sum_{a \in \delta(l_2)} z_a \geq 3 \quad \forall (l_1, l_2) \in \mathcal{I} \quad \text{WN} \quad (8)$$

Problem

- MST: k-BBD and WN constraints can not be integrated
- Generic solver: exponential number of constraints
- Polynomial algorithm: intractable [Gómez-Rodríguez et al. 2009]

Solving the ILP

⇒ Lagrangian Relaxation applied on constraints (7)-(8)

Lagrangian Dual Problem

$$\min_{u \geq 0} \max_{z \in Z} L(z, u)$$

Efficient minimization of the dual

- Min: Subgradient descent
- Max: Maximum Spanning Tree
- Many relaxed constraints: Non Delayed Relax-and-Cut

Efficient maximization of the primal

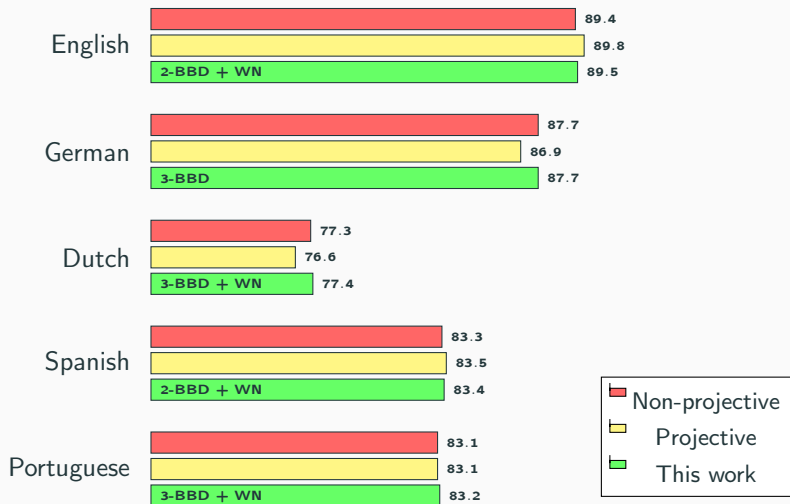
- Branch-and-Bound
- Problem reduction (exact pruning technique)

Distribution of dependency tree characteristics

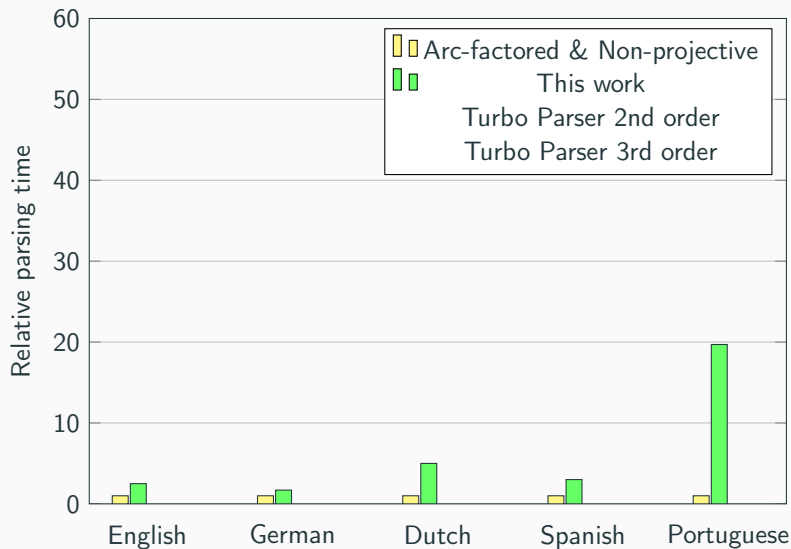
	English (PTB/LTH)		German (SPMRL)		Dutch (UD)	
	WN	WN	WN	WN	WN	WN
BD 1	92.26		67.60		69.13	
BD 2	7.58	0.12	27.12	0.79	28.50	0.08
BD 3	0.12	0.01	3.86	0.30	2.24	0.01
BD 4	0.00	0.00	0.19	<0.01	0.04	0.00
BD > 4	0.00	0.00	0.11	<0.01	0.00	0.00
	Spanish (UD)		Portuguese (UD)			
	WN	WN	WN	WN		
BD 1	93.95		81.56			
BD 2	5.99	0.04	13.92	0.05		
BD 3	0.02	0.00	3.76	0.02		
BD 4	0.00	0.00	0.54	0.00		
BD > 4	0.00	0.00	0.14	0.00		

- **Blue**: Projective dependency trees
- **Blue** + **Purple**: $\approx 99\%$ of the dependency trees

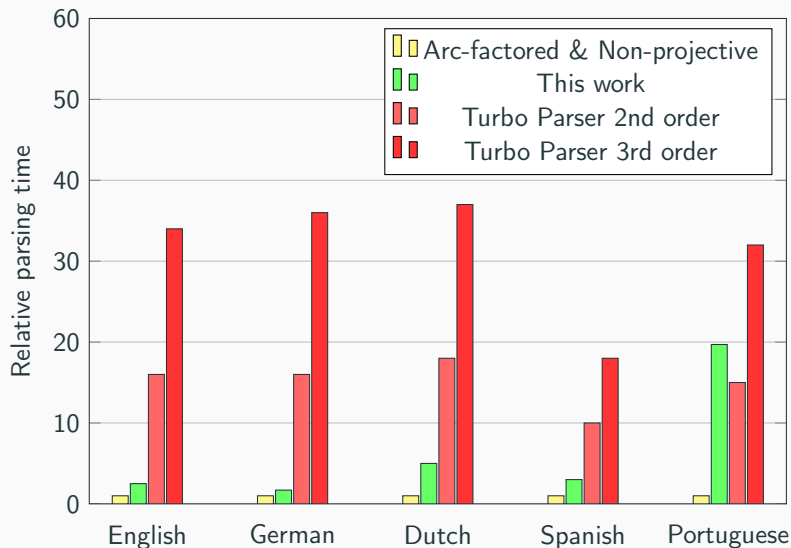
UAS (Ratio of correct arcs)



Efficiency: Relative parsing time



Efficiency: Relative parsing time



Conclusion: k-BBD and WN dependency parsing

Our contribution

- Novel characterization based on arc sets only
- The first efficient and flexible algorithm:
 - k-BBD with arbitrary k
 - WN optional } Tunable for different languages/properties
- First experimental results with K-BBD and WN parsing

Surprising observation

- Does not improve UAS under an arc-factored model

Perspectives

- LTAG derivation parsing (2-BBD and WN)
- Parsing lexicalized mildly context sensitive languages



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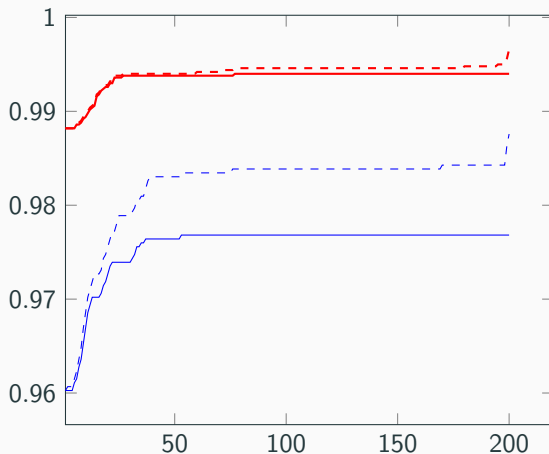


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Lagrangian Relaxation: Optimality Rate



(y-axis) Optimality rate

(blue) English

(solid) With certificate

(x-axis) Number of iterations

(red) German

(dashed) Without